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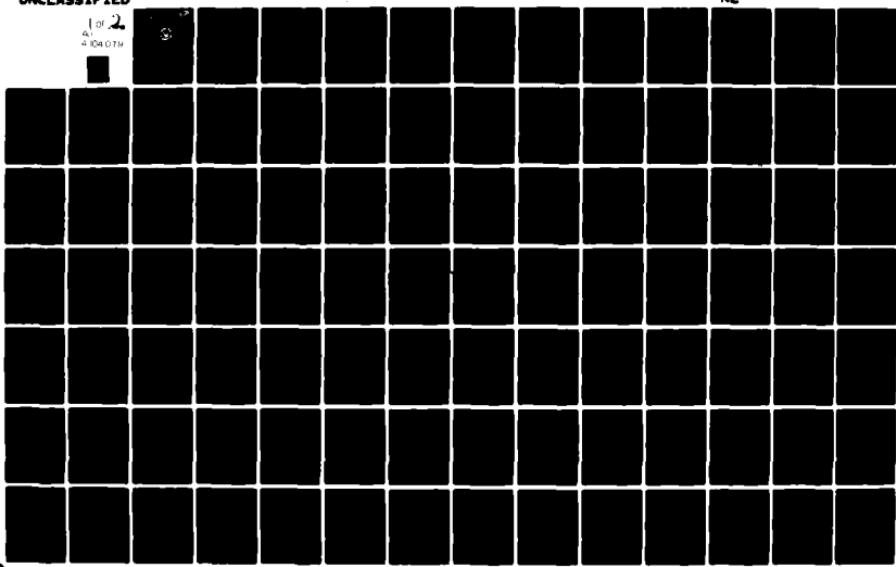
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THESIS

ALTERATION OF THE CP/M-86 OPERATING SYSTEM

by

Michael Bruno / Candalor

June 1981

1981

Thesis Advisor:

U. R. Kodres

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Alteration of the CP/M-86 Operating System

by

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Submitted in partial fulfillment of the  
requirements for the degree of

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from the

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June 1981

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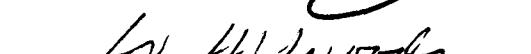
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## ABSTRACT

CP/M-86 is a microcomputer (INTEL 8086) operating system developed and marketed by Digital Research. The operating system is designed so that a user can adapt the system to his own input/output hardware devices. This thesis develops interfaces to two floppy disk controllers, the ISBC 281 (single density) and the ISBC 282 (double density) controllers. The interface includes the writing of a boot loader embedded in the ISBC 957 Execution Vehicle Monitor, the monitor system for the INTEL ISBC 86/12 single board computer. Also included is an interface module for the cold start loader (loader BIOS) and an input and output interface, BIOS. A design for the interface module of typical systems based on Winchester technology hard disks is also presented.

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## I. INTRODUCTION

### A. PURPOSE OF THIS THESIS

The adaptation of CP/M-86 to the hardware described herein was undertaken to provide an operating system for 8286 processor based single board computers at the Naval Postgraduate School. This operating system will support software development and system emulation for the AEGIS modeling project. The software will be available for general use at NPS. In addition the experience of modifying an operating system provided the author with an opportunity to learn about microcomputer hardware and microcomputer operating systems.

### B. HISTORY OF MICROCOMPUTER OPERATING SYSTEMS

This is a brief overview of the history of microcomputer operating systems summarized from Ref. 1. It is necessarily brief as the advent of microcomputer operating systems is itself rather recent. Microcomputers came of age with the construction of the entire central processor on one chip, the replacement of core memory with inexpensive mass produced semiconductor memory, the availability of the floppy disk and the standardization of diskette format. At first, the primary applications of microcomputers were in real-time control systems such as machine controlled tools. In such applications, process management is the main thrust

and system I/O is negligible. This required a simple, customized operating system. The first microcomputer operating systems, more properly called executive systems, were for real time applications. As microcomputer systems became less expensive, it became possible to devote a system to a single user as a program development tool. This use presented the need for higher level language support, which meant that an operating system had to interface one or more programming language(s) to the hardware. Several microcomputer manufacturers have produced their own operating systems. These operating systems are specifically designed for a "computer system" and are generally not user configurable.

Unlike the large, powerful operating systems found in mainframe and large minicomputer timesharing systems, microcomputer operating systems are relatively austere and simple. One of the primary reasons for this difference is that a microcomputer is usually a single user system (with some exceptions). As a result, the operating system does not need to provide features such as memory protection, process scheduling and time sharing of the CPU(s). Besides the simpler interface required of a microcomputer operating system, the operating system and the applications programs must function in a small amount of primary storage, typically between 16K and 64K, as compared to several megabytes in the large mainframes. Even though relatively

small and simple, a microcomputer operating system must still provide file management, process management and I/O management.

Two representative microcomputer operating systems are INTEL's ISIS-II and Digital Research's CP/M-80. To operate under ISIS, the user requires a minimum of 32K of primary storage. The CP/M user requires a minimum of 16K. Both provide the basic functions required of an operating system. ISIS, however, will only run on an INTEL computer system configuration and is not user modifiable. CP/M-80 is designed to run on any 8080 or Z-80 based microcomputer system after the user has modified the program module containing the hardware dependencies. This factor alone makes CP/M popular and has resulted in the production of many CP/M compatible utility and application programs by other companies. ISIS has some features beyond those of CP/M in the area of development software for INTEL hardware. CP/M's dynamic debugger (DDT), however, is more powerful and easier to use than INTEL's ICE system. Both ISIS and CP/M support essentially the same file operations. Currently, because of its flexibility, CP/M is the most widely used microcomputer operating system.

Multi-user systems such as MP/M and microcomputer network systems such as CP/NET (both produced by Digital Research), are now available.

### C. ADAPTATION TO THE USER'S ENVIRONMENT

Digital Research has attempted to make their CP/M operating systems as flexible, in terms of hardware suite, as possible. The method used is modular programming. The user interface, the Console Command Processor (CCP) has no hardware dependencies other than the CPU. The file management system, the Basic Disk Operating System (BDOS), is also independent of hardware. Both the CCP and the BDOS are interfaced to the Basic Input/Output System (BIOS) through logical I/O devices and logical disk devices. The BIOS, then, contains the logical device to physical device translation routines. Adaptation of the operating system to a unique environment requires only the modification of the appropriate BIOS routines, greatly simplifying the alteration process.

Once one has successfully completed one adaptation, follow-on adaptations will be much easier to achieve as an understanding of the operating system and its interface procedures is developed along with a better understanding of microcomputer architecture in general.

### D. ORGANIZATION OF THIS THESIS

This thesis is organized as a blueprint for alteration of the CP/M-86 operating system to any specific hardware configuration. This methodology will also serve, at least in general, for the alteration of any operating

system-to-hardware interface. Chapter 1 is a brief introduction to microcomputer operating systems in general and the modification of the CP/M-86 operating system in particular. Chapter 2 reflects the investigation of the candidate operating system in order to understand how to adapt it to the existing hardware. Chapter 3 is a summary of the study of the typical floppy disk or Winchester technology disk and a look at possible hardware candidates. Chapter 4 covers the adaptation of the I/O interface module (BIOS) and the bootstrap program for these versions of the operating system. Chapter 5 discusses some of the difficulties encountered and a plan for adapting CP/M-86 to a hard disk. The appendices contain the programs developed as part of thesis and one of the programs which was used as a model.

## II. STRUCTURE OF CP/M-86

### A. OVERVIEW

CP/M-86 is a microcomputer operating system for INTEL CORPORATION'S 8086 processor based microcomputers. It is the logical successor to CP/M-80, a similar operating system developed and marketed by Digital Research for the INTEL 8080 processor. File compatibility has been preserved with all previous versions of CP/M. CP/M provides a general environment for program construction, storage, editing, execution and debugging. The file structure of version 2 of CP/M-80 is used, allowing as many as sixteen drives with up to eight bytes on each drive.

CP/M-86 offers built-in utility commands, system transient commands and the capability of executing user defined transient commands (programs). Among the system transient programs are an Intel compatible assembler (ASM86) and a dynamic machine language program debugger (DDT). They are described in detail in Digital Research's publications [Ref. 2] and [Ref. 3] respectively.

A powerful feature of CP/M is its modularity. One of the three modules of the operating system, the Basic I/O System (BIOS), defines the hardware environment for the system. As a result of this modularity, CP/M-86 can be modified to run on any 8086/8088 processor based, single processor computer

system by merely changing the BIOS. A more detailed description of CP/M and its features is contained in Digital Research's publications [Ref. 4], [Ref. 5] and [Ref. 6].

### B. ORGANIZATION OF CP/M-86

The sources of CP/M-86 information for this paper are [Ref. 4], [Ref. 5] and [Ref. 6]. This chapter briefly summarizes the relevant material to this thesis.

The operating system is contained in file "CPM.SYS". "CPM.SYS" contains three program modules: the Console Command Processor (CCP), the Basic Disk Operating System (BDOS), and the user-configurable Basic Input/Output System (BIOS). This modularity allows the CCP and BDOS to be independent of the hardware in which the system is implemented.

The CCP is the system's interface to the user's console. It translates the user's commands into CP/M system calls in order to carry out the desired action. The BDOS module provides all the disk and file management. The BIOS contains all the hardware dependent features and interfaces. The operating system executes in any portion of memory above the interrupt locations, while the remainder of the address space is partitioned into as many as eight non-contiguous regions, as defined in a table in the BIOS.

CP/M-86 is too large a program to fit in the first two (system) tracks of a diskette. As a result the boot loader

loads into memory a cold start loader, called "LOADER.CMD", from the first two tracks. The boot loader makes the appropriate initializations and then transfers program control to the cold start loader. The cold start loader, which is essentially a subset of "CPM.SYS", finds "CPM.SYS" on the system disk, loads it into memory, makes the proper initializations, and finally transfers control to the operating system.

#### C. CCP BUILT-IN & TRANSIENT COMMANDS

The operation of CP/M-86 is similar to that of CP/M-80. Upon cold start the operating system signs on and drive A is logged-in, CP/M-86 then waits for an input command line. There are five built in commands:

- DIR - displays the directory of the designated drive
- ERA - erases the specified directory entry on the designated drive
- REN - renames the designated file
- TYPE - types the designated file to the logical console device
- USER - changes user directories in multi-directory systems

Also the command line may begin with the name of a transient program with the assumed file type of CMD. CMD stands for "command file" and is used to differentiate CP/M-86 transient command files from COM files under CP/M-80

which serve the same purpose. Transient programs are loaded into memory in the Transient Program Area(s) (TPA), as defined in the BIOS, in stack order.

CP/M-86 supports programs written in three memory models: the 8080 model, the Small model and the Compact model.

The 8080 model supports programs which are directly translated from CP/M-80 when code and data areas are intermixed. The model consists only of a code group which, in turn, is normally a single segment of 64K or less. The operating system and the cold start loader are written in the 8080 model.

The Small model supports programs where there is a separate code and data group. Normally the Small model programs are 64K or less.

The Compact model occurs when any of the extra, stack or auxiliary groups are present in the program. Each group may consist of one or more segments.

The three models differ primarily in the manner in which the segment registers are initialized upon transient program loading. The operating system's program load function determines the memory model used by the transient program by examining the program group used. All three models are discussed in more detail in the next section.

### 1. Transient Program Execution Models

The initial values of the segment registers are determined by the "memory model" of the transient program and are described in the CMD file header generated by the program "GENCMD.CMD" or "GENCMD.COM". The three models are depicted in Figure 1.

```
-----  
! 8080 Model ! Code and Data Groups Overlap      '  
-----  
! Small Model ! Independent Code & Data Groups    !  
-----  
! Compact Model ! Three or More Independent Groups !  
-----
```

Figure 1 Transient Program Memory Models.

#### a. The 8080 Model

The 8080 Model is assumed when the transient program contains only a code group (containing both code and data). In such cases, the CS, DS and ES registers are all initialized to the beginning of the code group, while the SS and SP registers remain set to a 96-byte stack area in the CCP. The Instruction Pointer (IP) is set to 100H, similar to CP/M-80. The intermixed code and data regions are indistinguishable. This model allows simple translation of 8080, 8085 and Z80 code into the 8086 and 8088 environment. Following program load, the 8080 Model appears as in Figure 2, where low addresses are shown at the top of the diagram.

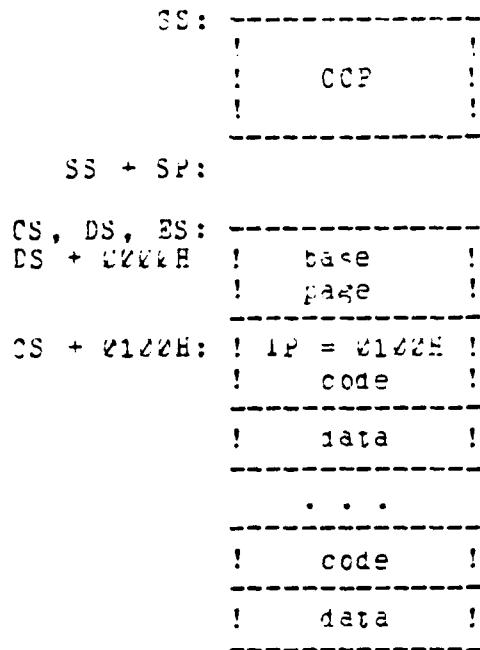


Figure 2 The 8086 Memory Model.

### b. The Small Model

The Small Model is assumed when the transient program uses both a code and data group. (In ASMB6, all code is generated following a CSEG directive, while data is defined following a DSEG directive.) In this case CS is set to the beginning of the code group, the DS and ES registers are set to the start of the data group, and the SS and SP registers remain in the CCP's area as shown graphically in Figure 3.

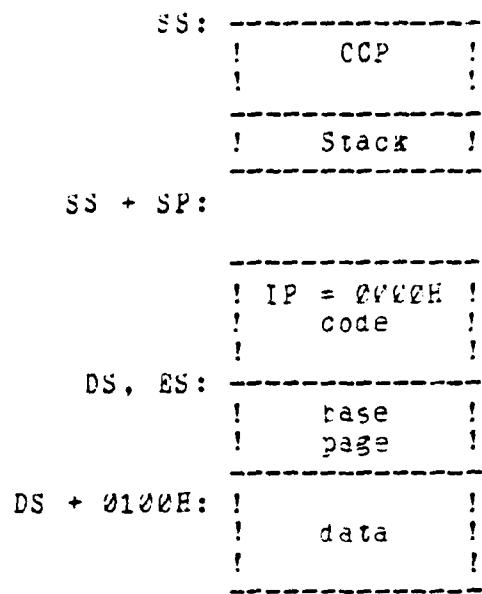


Figure 3 The Small Memory Model.

#### c. The Compact Model

The Compact Model is assumed when separate code and data groups are present, along with one or more of the remaining groups. In this case, the CS, DS and ES registers are initialized to the base address of their respective areas. The SS and SP registers remain in the CCP area. If the user intends to use the stack group as a stack area, the transient program must set the SS and SP registers upon entry. The initial configuration of the segment registers in this model is shown in Figure 4.

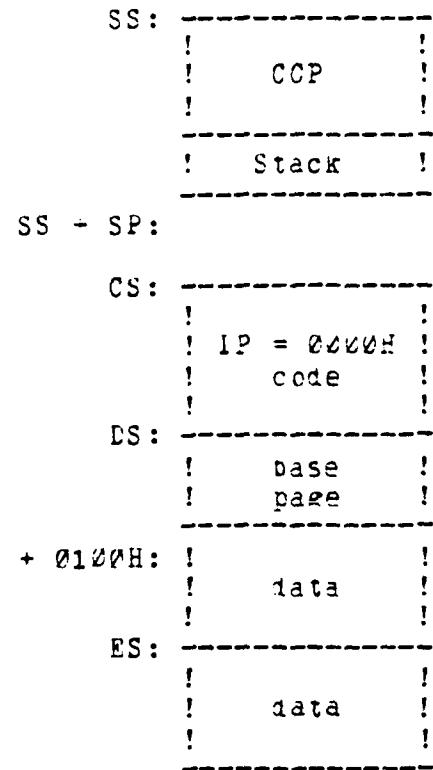


Figure 4 The Compact Memory Model.

The values of the various segment registers can be programmatically changed during execution by changing the values in the base page as described in the preliminary documentation, thus allowing access to the entire memory space.

## 2. Transient Program Setup And Termination

Similar to CP/M-80, the CCP parses up to two file names following the command and places the properly formatted File Control Blocks (FCB's) at locations 2005CH and 006CH in the base page relative to the DS register. Under CP/M-80, the default DMA (direct memory access)

address is initialized to 0080H in the base page. Due to the segmented memory of the 8086 and 8088 processors, the DMA address is divided into two parts: the DMA segment address and the DMA offset. Also, under CP/M-86, the default DMA base is set to the value of DS, and the default DMA offset is initialized to 0080H. Thus, CP/M-86 and CP/M-88 operate in the same way in that they both assume the default DMA address is the second half of the base page.

The CCP transfers control to the transient program through an 8086 "Far Call." In all but one case of the Compact Model, the transient program may choose to use the 96-byte CCP stack, and optionally return directly to the CCP upon program termination by executing a "Far Return." Programmatic termination also occurs when BDOS function zero is executed. The operator may terminate program execution by typing a single CONTROL-C during line edited input. This has the same effect as programmatic execution of BDOS function zero. Contrary to the operation of CP/M-80, no disk reset occurs and the CCP and BDOS modules are not reloaded from the disk upon program termination. In short, for the user familiar with CP/M-80, the CP/M-86 environment is very similar, but more powerful.

#### D. BDOS SUMMARY

Entry into the BDOS is made through the 8086 software interrupt # 224. The BDOS is, essentially, a set of 59

functions of three basic types; simple functions, file operations and extended operations. The interface convention for BDOS calls requires that function code be passed in register CL with parameters passed in register DL or DX depending on whether it is a byte or word value. Byte values are returned in the AL register and word values in registers AX and BX. Table 1 below, from Reference 6, summarizes these conventions. A full description of each BDOS function is given in [Ref.6].

! BDOS Entry Registers !		! BDOS Return Registers !	
! CX Function Code	!	! AL Byte Value	!
! DL Byte Parameter	!	! AX Word Value	!
! DX Word Parameter	!	! BX Word Value	!
! DS Data Segment	!	! ES Double Word Offset	!
!	!	! FS Segment Address	!

Table 1 BDOS Parameter Conventions.

## E. BIOS SUMMARY

The BIOS is loaded into memory just above the CCP and BDOS modules as illustrated in Figure 5.

Since the BIOS may be configured by the user, it may vary somewhat in length. Individual routines within the BIOS may be at different memory locations. In order to standardize the interface to the BIOS, all accesses to the BIOS are made through the jump vector at the beginning of that module. The BIOS, like the BDOS, also has parameter

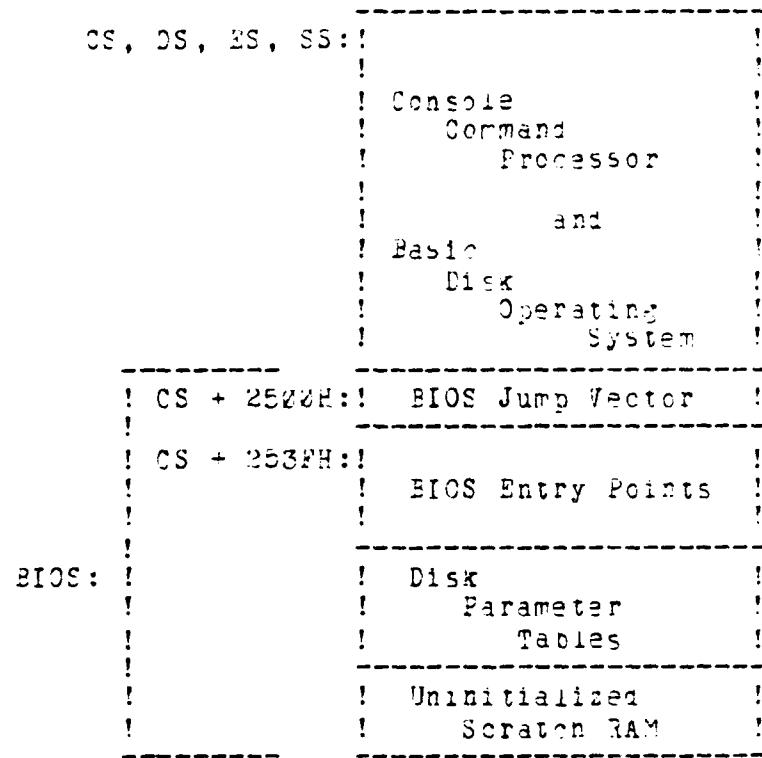


Figure 5 Memory Location of the BIOS.

passing conventions. Parameters for routines in the BIOS are passed in the CX register and the DX register when required. Byte values are returned in the SI register and word values in BX.

There are three major types of routines in the BIOS: system initialization/reinitialization, simple character I/O and disk I/O. All simple character I/O operations are assumed to be in ASCII, both upper and lower case, with the high order (parity) bit set to zero. CP/M sees all peripheral devices as "logical" devices. Translation from logical device selection to physical device assignment is

accomplished in the BIOS, thus isolating the CCP and BDOS from hardware dependencies. BIOS routine entry is explained in Digital Research's publication [Ref. 5]. The BIOS also contains the Disk Parameter Tables which contain the description of the disk drive and provide a scratchpad area for certain BDOS operations.

### III. INPUT/OUTPUT DEVICES

In CP/M-86 the CCP and BDOS accomplish all I/O via four "logical" devices. The BIOS assigns whatever physical devices are in that particular system to those logical devices. This mapping in the BIOS preserves the independence of the CCP and BDOS from the hardware configuration.

#### A. LOGICAL I/O DEVICES

CP/M-86 addresses four logical I/O devices: the console, the list device, the punch device and the reader. The console is the principal interactive peripheral through which the operating system communicates with the operator. The list device is the principal listing device, usually a hardcopy printer. The punch device is the principal tape punching device, usually a high-speed paper tape punch or teletype. The reader is the principal tape reading device. When the "IOBYTE" function is implemented, dynamic logical to physical device mapping may be accomplished as described in Ref. 6.

#### B. PHYSICAL I/O DEVICES

The CONIN, CONOUT, LISTOUT, PUNCH and READER routines in the BIOS define the physical interfaces with peripherals. The system adapter may define, in the BIOS, such devices as cassette tape recorders etc. so long as it is interfaced

with one of the logical devices. In this adaptation the list device and the console device are both mapped to the serial eige connector where the CRT console is connected. The reader is "stubbed" with an "end of file" input, that is, instead of a routine to interface a physical read device, the BIOS simply returns an indication that the read has been complete. And the punch device map is "stubbed" with a return statement.

## C. DISK DEVICES

### 1. Hard Disks, floppy Disks

There are many implementations of the hard disk technologies. There are fixed and movable head disks, removable disk packs and even combination hard and floppy systems. Floppy diskettes come mainly in the 5" and 8" size, single and double density, single and double sided, and as indicated above in combination with hard disks.

### 2. Organization of Data

Although each disk drive may be different, data is stored in conceptually the same manner. The disk surface is divided into tracks (or cylinders, if a multi-platter system.) Each track is divided into sectors. Each sector is addressable by the controller, making it the basic unit of storage. In multi-platter and/or multi-head systems, to access the disk the controller must select the proper head/platter as well as the track and sector required.

The amount of data that can be stored on a device is dependent on the size of the device and the recording format. Double density, as the name implies, gives twice as much storage on a diskette as single density. The cost, however, is greater.

Although the basic unit of storage is the sector, sectors are not the same size in every system. In general, the larger the sector, the more efficient the storage, but the less efficient the access. Many systems allow the user to select the sector size from a limited set of choices. Sectors are normally a multiple of 128 bytes.

### 3. Interfaces to the Computer

The key to the storage of information on the recording media, at least from the operating system modifier's point of view, is the disk drive controller. The controller itself is usually a microprogrammed microprocessor. The controller handles the actual reading from and writing to the disk in addition to other functions such as seek, format etc. The relative autonomy of the controller frees the operating system from having to handle disk I/O on a primitive level. However, the BIOS, which is hardware specific, must still communicate with the controller at a fairly low level.

Most microcomputer system I/O is done by DMA. In general the host operating system creates, somewhere in memory, an entity, often called a "command packet" or "I/O

parameter block" or some similarly descriptive name. The "packet" is usually seven to ten bytes of information which contain the detailed command for the disk drive controller. These "packets" form the sole means of issuing I/O commands to the controller.

Normally the disk drive controller/interface shares a bus with the host system. As a result the controller's command/status registers have device addresses from the bus. In most systems, they can be set by the user prior to system start-up.

The host system sends the address of the I/O command packet to the command registers of the controller. Upon receipt of this address the controller initiates action to gain control of the bus. When the controller has control of the bus it reads the appropriate number of bytes from the address it was given. The controller decodes this information and then carries out the prescribed operation. The controller may signal completion in various ways, the most common being entering a completion code in the command packet for the host to read, sending an interrupt to the host processor, or storing the status in an on-board status register for the host to read.

Many systems allow the DMA to be "throttled", that is, the controller gives up control of the bus periodically in order to increase overall system speed.

Other features commonly included in disk drive controllers are: linked I/O, that is, the ability to execute more than one I/O command packet without prompting from the host processor. Multiple sector I/O, that is, the ability to read or write more than one sector in response to a single I/O command packet.

#### 4. Examples of Particular Controllers

##### a. iSBC 201 (Single Density MDS)

The iSBC 201, as described in Ref. 7, is the controller/interface for INTEL's INTELLEC MDS 808, an 8080 processor based microcomputer development system.

(1) iSBC 201 Controller Operation. The controller is composed of two circuit boards, a channel board and an interface board. They interface with the host processor via the system MULTIBUS, a system's bus used by INTEL Corporation. The channel board and interface board together handle all communications between the host CPU and the diskette system. They contain an 8-bit microprogrammed processor which can access system memory for obtaining channel commands via DMA. The controller also monitors the disk subsystem status and error conditions and makes their status available to the host CPU.

This diskette system records data by the Frequency Modulation (FM) method, giving a formatted 8" diskette capacity of approximately 256K bytes, divided into 77 tracks of 25 sectors each.

Functionally, the host CPU must create a command packet in memory for each operation. INTEL calls this packet an I/O Parameter Block (IOPB). An IOPB is ten bytes in length and specifies all the details of the diskette operation to be performed. The CPU, in the case of CP/M-86, through the BIOS module, sends the address of the IOPB to the controller. Then the controller gains control of the bus, retrieves the IOPB and executes the command. Upon completion the controller posts the diskette subsystem status and, if enabled by the IOPB, sends a completion interrupt to the host CPU. The information in the IOPB consists of:

- Byte 1 - the channel word, this byte specifies the enabling of the lock override, random format of the lock override, random format sequence, interrupt control, data word length, successor bit, branch on wait and wait bits.
- Byte 2 - specifies the drive selected, data length (8 or 16 bits/word) and the operation to be performed.
- Byte 3 - specifies the number of sectors to be transferred.
- Byte 4 - specifies the target track number (0-77).
- Byte 5 - specifies the first sector to be accessed (1-26).

Byte 6 - specifies the least significant byte of the buffer address.

Byte 7 - specifies the most significant byte of the buffer address.

Byte 8 - indicates a block number which allows a unique identification of an IOPB during linked IOPB operations.

Byte 9 - contains the least significant byte of the buffer address of the next linked IOPB.

Byte 10 - contains the most significant byte of the buffer address of the next linked IOPB.

The ISBC 201 can execute seven commands:

- 1) recalibrate (seek track 0)
- 2) seek
- 3) format a track
- 4) write data (without address marks)
- 5) write data
- 6) read data
- 7) verify CRC

The controller has seven registers that are accessible to the host CPU. The host CPU can read three of the registers: The Result Status register indicates the status of both drives (ready or not ready), the status of the controller for that drive (present or not present), and the status of the controller's interrupt flip-flop

(interrupt pending or completed). The Result Type register indicates whether the Result Byte register contains I/O error codes or ready status. The Result Byte holds the I/O error codes or diskette drive status. The host CPU can write to four of the controller's registers: Writing anything to the Reset Diskette System register resets the entire diskette subsystem. Writing to the Stop Diskette Operation register terminates I/O after completion of the current operation. The Memory Address Lower register receives the least significant byte of the address of the IOPB. The Memory Address Upper register receives the most significant byte of the IOPB address and when written into signals the controller to retrieve the IOPB and commence the specified operation.

(2) BIOS Use of the iSEC 201. The CP/M-86 BIOS uses only operations 1, 5 and 6 (seek is implicit in read and write operations). In addition, CP/M-86 does not use linked IOPB's and only does single sector disk accesses. This very much simplifies the I/O routines in the BDOS and the BIOS. Not using the linked IOPB capability allows reducing the IOPBs to the first seven bytes, of which bytes 1 and 3 remain constant. Byte One remains unchanged because the mode of disk access remains unchanged. Byte Three, the number of sectors, remains set at one, and the operating system is freed from computing the number of sectors per

access. These simplifications allow the BIOS to have a single IOPB template in memory.

A limitation of the iSBC 201 is its 16-bit addressing. This limitation means that the controller can only address 54K of system memory as compared to the 8086 processor's megabyte of address space. As a result, the external address of the iSBC 86/12 must reside in the first 54K of the megabyte (from 20000H to 6FFFFH). The BIOS in this adaptation converts the segment and offset address provided by the BDOS into a 16-bit physical address for the controller.

(3) Bootstrap Use of the iSBC 201. The bootstrap program does use the multi-sector access capability of the controller for loading the cold start loader. This requires four IOPBs in the bootstrap program but reduces the number of disk accesses from 53 to four. Considering the specialized function of the bootstrap loader and its lack of interface with the BDOS, this is a very efficient deviation from the otherwise efficient CP/M method of disk access.

#### b. iSBC 202 (Double Density MDS)

The iSBC 202 is the controller/interface for INTEL's INTELLEC MDS 888 microcomputer development system. It is described fully in Ref. 8.

(1) iSBC 202 Controller Operation. From the users point of view this controller is essentially the same as the iSBC 201. The main difference is the recording

format. Modified-Modified Frequency Modulation (MMFM) is used, allowing the same media to hold (formatted) 512K bytes of data, divided into 77 tracks of 52 sectors each. This is twice the capacity of the single density system.

(2) BIOS Use of the iSSC 222. The interface to the controller is the same as that of the iSSC 221. The difference in organization and capacity is only evident in the disk definition table "DOUBLE.LIB".

(3) Bootstrap Use of the iSSC 222. CP/M's double density formatter formats the first two tracks of a diskette in single density, ie. 26 sectors per track. The cold start loader fits in the first two tracks of a double density in the same way as in single density. As a result, the save bootstrap program will load the cold start loader from both single and double density diskettes.

#### c. REMEX RDW 3200

The RDW 3200, as described by Ref. 9, Ref. 10 and Ref. 11, is a multi drive unit consisting of a fixed Winchester Technology 14" disk and two 5" flexible diskette drives. The diskette drives are "jumper" selected as either single or double density. In both types the sector size is selectable. The formatted capacity of the fixed disk with sector size set at 128 bytes is 10 megabytes. This data is on 210 tracks of 124 sectors for each of two read/write heads. The single density floppy drives, formatted for 128 bytes per sector, hold 26 sectors on each of 77 tracks for a

total of 256K bytes of storage. Set for double density, the smallest sector size available is 256 bytes. At 26 sectors per track, for 77 tracks, formatted storage is 512K bytes. If this drive were used for CP/M-86 in the double density mode, the difference between diskette sector size (256 bytes) and CP/M-86 sector size (128 bytes) would be handled by a "blocking/deblocking" algorithm like the one provided with CP/M-86.

(1) The RDW Controller. The heart of the controller is a microprogrammed Motorola 6826 8-bit microprocessor. The controller physically resides inside the RDW frame and is linked to the host system by an interface card. This alteration utilized a MULTIBUS interface, which resided in the host's system MULTIBUS. The interface provides registers for communication between the host and the controller CPU's. Data can be handled as 8-bit words, 16-bit words or as 8-bit half-words. The controller can accomplish I/O by DMA, programmed I/O or by interrupts. All disk writes are by Modified-Modified Frequency Modulation (MMFM). The disk drive system can also be DMA throttled, which permits other masters to gain access to the system's bus in between accesses by the disk unit.

Functionally, the host CPU must create a command packet in memory for each operation. A command packet is six to fourteen bytes in length and specifies all the details of the disk operation to be performed. In the

DMA mode the host CPU must test the status register in the controller interface to assure that the controller is ready. When the controller is ready the CPU, in the case of CP/M-86, through the BIOS module, sends the address of the command packet to the controller interface. Then the controller gains control of the bus, retrieves the command packet and executes the command. Upon completion, the controller posts the disk subsystem status in the command packet in system memory and, if enabled by the command packet, sends a completion interrupt to the host CPU. The command packet consists of six to fourteen bytes. This controller supports five types of operations. The size of the packet and the information it contains are determined by the operation to be performed. The five operations supported are:

- 1) read data/write data
- 2) write I.D. and data for single record  
(fixed disk only)
- 3) copy from one drive to another
- 4) format designated disk
- 5) maintenance package

The controller has four registers that are accessible to the host CPU. The base address of these registers is switch selectable. The base address plus one is the status register, from which the host CPU determines

system status. The base address plus three receives the lower byte of the address of the command packet. The base address plus two receives the middle byte of the command packet address. The base address receives the upper byte of the packet address (RDW 3200 supports 24-bit addressing) and when written into signals the controller to start DMA.

(2) BIOS Use of the RDW 3200. The CP/M-86 BIOS would use only the read/write operation. The fact that the hard disk has more than one head would require that the BIOS disk definition table look like one continuous set of tracks and that prior to initiating DMA, the BIOS translate a logical track number to a physical head and track number. The read and write packets have the same format which requires only one packet template in the BIOS. That packet takes the following form; indicated as 16-bit words:

Word 0 - I/O modifiers (linked I/O, interrupts, etc.),  
operation and drive selected.

Word 1 - status word - written by controller.

Word 2 - track number.

Word 3 - head and sector start number.

Word 4 - lower 16 bits of DMA address.

Word 5 - high byte of DMA address.

Word 6 - transfer word count.

Although the RDW supports 24-bit addressing, it requires a 24-bit physical address, not the

segment and offset type address provided by the BIOS. Therefore the BIOS must translate the addresses before placing them in the command packet and before sending them to the interface.

(3) Bootstrap Use of the RD<sub>n</sub> 3222. The bootstrap program would use the multi-sector access capability of the controller for loading the cold start loader (the command packet specifies the number of words to be transferred). If the operating system were to be loaded from a diskette, the bootstrap operation would be very much like that described for the iSBC 201. For a system load from the hard disk the bootstrap program could load the operating system without the use of a cold start loader. This would only require two disk accesses, one to determine the load location and the other to actually load "CPM.SYS".

#### IV. ALTERATION OF CP/M-86

##### A. CHANGES REQUIRED TO IMPLEMENT CP/M-86

As distributed, CP/M-86 is set up for operation with an Intel SBC 86/12 microcomputer and an Intel SBC 274 diskette controller with a Shugart SA-200 floppy disk drive. Since CP/M-86 is modular, only the BIOS need be modified for "non standard" hardware. The distribution version includes source code for its BIOS and a skeletal BIOS to aid in the construction of a customized version. Although the distribution version does not provide a bootstrap ROM, the source code for the program is provided. This source code provided an example for the creation of a customized bootstrap program. The bootstrap ROM is available from Digital Research.

The changes required to customize the BIOS can be divided into four types. The first consideration is the computer selected for the implementation. If an 8086/8088 based computer other than the iSBC 86/12 were chosen, the computer initialization, including the constant definitions for USART ports and character I/O routines such as console status, console input and console output, have to be changed to match the host hardware. Since the iSBC 86/12 was used, no changes were required in this portion of the BIOS. Second, if the disk drive controller or other DMA device is

not an iSBC 204, the controller port definitions and the routines which actually communicate with the controller must be altered. The "execute" and "sendcom" routines were the bulk of the modification. These routines check system status, translate system commands to the language of the controller, deliver the commands to the hardware and handle any hardware errors. Third, if any other serial or parallel I/O device is to be used, the appropriate initialization and execution routines must be written. The fourth consideration is the disk definition table which is assembled with the BIOS via an "include" statement. Disk parameter tables must be created to describe the disk system. Disk parameter tables are discussed in the next section. In this version only the second and fourth types of modifications were necessary and those changes are reflected in Appendix A (single density) and Appendix B (double density). Appendix D contains the distribution BIOS. After assembling the BIOS, the hexadecimal code, "BIOS.H86", is appended to "CPM.H86" and a command file is generated by the method described in Ref. 6 using the GENCMD utility. The file created is named "CPM.SYS" and is the operating system.

#### B. DISK PARAMETER TABLES

The disk parameter table serves to define the organization of the storage media for the BDOS file management functions. The disk definition consists of the

sequence of statements in Figure 6 (as shown in Ref. 6). The DISKS statement defines the number of drives in the system, with n being an integer from 1 to 16. A series of DISKDEF statements follow. Each statement defines the characteristics of a logical disk, 0 through n-1. DISKDEF statements are formed as defined in Ref. 6. The format is shown in Figure 7.

```
DISKS    n
DISKDEF  0, ...
DISKDEF  1, ...
.....
DISKDEF  n-1
.....
ENDEF
```

Figure 6 BIOS Disk Definition File.

DISKDEF dn,fsc,lsc,[skf],ols,dks,dir,cks,ofs,[v]

where

dn	is the logical disk number, 0 to n-1
fsc	is the first physical sector number (0 or 1)
lsc	is the last sector number
skf	is the optional skew factor
ols	is the data allocation block size
dks	is the disk size in ols units
dir	is the number of directory entries
cks	is the number of checked directory entries
ofs	is the track offset to logical track v
[v]	is an optional 1.4 compatibility flag

Figure 7 DISKDEF Statement Format.

The disk tables may be generated by hand or by executing the GENDEF utility program. The table provided with the

distribution version, called "SINGLES.LIB", was generated from the source file "SINGLES.DEF" by the GENDEF utility running under CP/M-86. This table was correct for the single density implementation. It was necessary to create a new table for the double density system. This file is called DOUBLE.DEF. Table generation is described fully in Section 6 of Ref. 5. The disk parameter tables are listed in the BIOS right after the "include" statement (see Appendices A and B).

### C. COLD START

#### 1. The Cold Start Loader

Since CP/M-86 is too large to fit in the first two (system) tracks of a diskette, it is loaded into memory in two steps. First, a cold start loader is loaded from the first two tracks into memory. Next the loader loads the operating system and transfers control to it. The loader ("LOADER.CMD") is a simplified version of CP/M-86 with enough power to locate the operating system file "CPM.SYS" on the current disk, make the proper initializations, load CP/M-86 into memory and then transfer program control to it. The loader is created from files LDCPM, IDBDOS and the loader version of the BIOS. The loader BIOS is generated from the same source code as the BIOS by setting the software switch "LOADER\_BIOS" equal to true prior to assembly.

The loader program is moved to the first two tracks of a diskette by the LDCCOPY utility if running on a working CP/M-86 system. If development is done on a CP/M-82 system this can be accomplished with the DDT and SYSGEN utilities. Ref. 6 errs in its description of the latter procedure. The correct procedure is described in the next chapter.

## 2. The Bootstrap ROM

In order to set the cold start loader into memory, there must be a bootstrap loader of some kind. This boot loader must initialize the programmable chips on the single board computer and the disk drive controller which will access the operating system disk. It then loads the first two tracks of the diskette in the system disk drive into memory and then transfers control to the program loader, "LOADER.CMD". The bootstrap program is normally resident in a read only memory (ROM) or electrically programmable ROM (EPROM) and is then referenced to as the boot ROM.

The distribution version of CP/M-86 also contains the listing for a bootstrap ROM (ROM.A86). The boot ROM itself is available from Digital Research. When installed, it becomes part of the 8086 address space. Upon system reset, the processor begins execution at effective address OFF000H, which is the top paragraph of the iSBC 86/12 EPROM space. The bootstrap program is hardware dependent which necessitated the creation of a customized initial loader for this implementation.

Intel's SBC 957 Execution Vehicle Monitor (EVM) occupies the EPROM locations when installed in the iSBC 86/12 and is currently in use at the Naval Postgraduate School. In order to retain the use of the iSBC 957 and to simplify implementation, the customized bootstrap program has been embedded in a free area of the EVM's EPROMs. Since the monitor initializes the single board computer when it is started, the CP/M-86 bootstrap task is simplified. The bootstrap program listing is in Appendix C. It is a modified version of the "debug" version of Digital Research's ROM program. The modified bootstrap program is located at effective address 0FFD40H. It may be executed from the EVM by executing the command GFFD4:0 or its equivalent.

## V. CONCLUSIONS AND RECOMMENDATIONS

### A. ADAPTATION DIFFICULTY

Modification of CP/M-80 is a straightforward simple procedure if one is familiar with CP/M on a system's software level and with at least some representative hardware. If one does not have such a background (the author did not), the task is not overwhelming, but considerably more difficult. The novice will probably invest much time and effort in investigating "dead ends" because of not understanding the logical design of the operating system. A particularly vexing problem encountered in the first adaptation was that in the later stages of development, every error in the corrected software seemed to destroy the information on the diskette, making debugging difficult and requiring frequent regeneration of software. During this period of "destructive testing" approximately 90% of the time and effort were spent on such overhead and only 10% on actual debugging. The real problem there was not the time lost but the interruption in the train of thought.

Documentation inadequacies are another source of problems. The alteration guide for CP/M-86 provided by Digital Research (Ref. 6) assumed a thorough knowledge of CP/M-80, which was not possessed by the author. The CP/M-86 documentation also seemed to assume a thorough knowledge of

the operating system's modules. In addition, there were several errors in the alteration guide.

The procedure for moving the cold start loader to tracks zero and one under CP/M-80 is incorrect and if followed the first 802H bytes of the program will be lost. A correct procedure is to load the cold start loader with DDT, move the program so that it starts at 900H, exit DDT and finally call the SYSGEN utility. A correct sequence of commands looks like this:

```
DDT LOADER.CMD  
m1100,1800,1900  
m00,1100,1200  
m400,a00,c00  
m100,400,900  
<CONTROL-C>  
SYSGEN  
<CR>  
B  
<CR>
```

The documentation for the 8086 assemblers, "ASM86.COM" and "ASM86.CMD" also contains errors. According to the user's manual, [Ref. 2], the "device switch" for the listing device is "P". The correct switch is "Y".

The technical manuals provided with the disk drives and controllers used rather ambiguous and non standardized

terms. This often required experimentation to determine what was really meant.

Resolution of the above difficulties, however, was a good learning experience for the author.

## B. RECOMMENDATIONS FOR FUTURE HARD DISK ADDITION

### 1. Discussion

Although there are several methods of accomplishing disk I/O, DMA seems to be the simplest to implement and debug. A future hard disk addition would greatly enhance CP/M-86's usefulness. In this vein, a hard disk/floppy disk combination would be ideal. The combination of hard and floppy disks would provide the speed and storage capacity on one hand (from the hard disk) and the ability for the user to keep copies of his files where he is assured of their security and integrity. However, inclusion of the iSBC 201 or 202 is not recommended. The limited addressing capability of these controllers would hinder overall system effectiveness and force the processor to operate in the bottom 64K of the address space. As a rule of thumb, if more than one device is to be added to the basic system, only one device should be added at a time.

### 2. Template for Adaptation

Given that a hard disk is to be installed in place of the diskette system, the following procedure should be followed:

First, the CP/M-86 BIOS should be studied in conjunction with the current hardware to see how the interface is currently accomplished. The system modifier must understand how the operating system interacts with hardware before creating his own interface. Second, the target hardware must be studied. The electronics are not important, but what the hardware does logically and how it communicates with the controller is paramount. In particular, the organization of data on a disk drive must be thoroughly understood. If the organization of data is selectable, the most efficient and straightforward organization must be chosen. If it is not selectable and not directly compatible with the BDOS, a "clocking/deblocking" or some other scheme must be considered. Third, a disk definition must be written to reflect the logical organization of the disk. If the logical organization of data does not match its physical organization, the executing routine in the BIOS would have to make the translation. For example, in a multi head disk system, the tracks would have to be numbered in the disk definition as though they were on the same platter (logical org.), the BIOS would select a sector and a "logical" track for I/O, but before sending the channel command the BIOS would have translate that "logical" track number to a head and track combination. Fourth, a template for the channel command should be placed in the BIOS with appropriate variable names to allow the BDOS to

provide as much information directly as possible. Fifth, write the "execute" routine. This routine, the bulk of the coding, must complete the channel command, prepare the disk for access, send the activating command, check completion status and handle hardware errors. This step requires a good knowledge of the target disk system and is very much dependent on the disk chosen. Sixth, once the revised BIOS is written, it must be assembled (in the loader version too, if booting from a floppy disk). The files "CPM.H86", "BIOS.H86" and "PAT2.H86" are combined into "CPMX.H86". This resulting file is converted to executable form by executing the command "GENCMD CPMX 8080[A40]" as described in Ref. 6. The resulting file is then renamed "CPM.SYS".

The bootstrap program will be very simple. It can be written to explicitly read the first sector of the disk, to determine the loading target address, and to read the following 76 (128 byte) sectors. Once the BIOS has been modified, the bootstrap program will be almost a trivial subset of that code.

## APPENDIX A

title 'Customized Basic I/O System'

```
;*****  
;*  
;* This Customized BIOS adapts CP/M-86 to *  
;* the following hardware configuration *  
;* Processor: iSBC 8612 *  
;* Controller: iSBC 221 *  
;* Memory model: 8080 *  
;*  
;* Programmer: M.B. Candalor *  
;* Revisions : *  
;*  
;*****  
  
true equ -1  
false equ not true  
cr equ 0dh ;carriage return  
lf equ 0ah ;line feed  
max_retries equ 10 ;for disk i/o, before perm error  
  
;*****  
;*  
;* Loader_bios is true if assembling the *  
;* LOADER BIOS, otherwise BIOS is for the *  
;* CPM.SYS file.  
;*  
;*****  
  
LOADER_BIOS EQU TRUE  
bios_int equ 224 ;reserved BDOS interrupt  
  
IF not loader_bios  
-----  
;  
bios_code equ 2500h  
ccp_offset equ 0000h  
bdos_ofst equ 4B06h ;BDOS entry point  
;  
-----  
ENDIF ;not loader_bios  
  
IF loader_bios  
-----  
;  
bios_code equ 1200h ;start of LDBIOS  
ccp_offset equ 0003h ;base of CPMLOADER  
bdos_ofst equ 4426h ;stripped BDOS entry
```

```

;-----|
;-----|
        ENDIF    ;loader_bios

csts     equ 21an          ;I8251 status port
ciata    equ 0d8n          ;      iata
;
;
;***** INTEL iSBC 201 Disk Controller Ports *****
;*
;* INTEL iSBC 201 Disk Controller Ports
;*
;***** *****
;

base     equ    078n
rtype   equ    base+1
rbyte   equ    base+3
reset   equ    base+7

dstat   equ    base
ilow    equ    base+1
inighn  equ    base+2

        cseg
        org    ccpoffset
ccp:
        org    bios_code

;***** *****
;*
;* BIOS Jump Vector for Individual Routines *
;*
;***** *****

        jmp INIT           ;Enter from BOOT ROM or LOADER
        jmp WBOOT          ;Arrive here from BDOS call 0
        jmp CONST           ;return console keyboard status
        jmp CONIN           ;return console keyboard char
        jmp CONOUT          ;write char to console device
        jmp LISTOUT         ;write character to list device
        jmp PUNCH           ;write character to punch device
        jmp READER          ;return char from reader device
        jmp HOME            ;move to trk 00 on cur sel drive
        jmp SELDSK          ;select disk for next rd/write
        jmp SETTRK          ;set track for next rd/write
        jmp SETSEC          ;set sector for next rd/write
        jmp SETDMA          ;set offset for user buff (DMA)
        jmp READ             ;read a 128 byte sector
        jmp WRITE            ;write a 128 byte sector
        jmp LISTST           ;return list status

```

```

jmp SECTRAN      ;xlate logical->physical sector
jmp SETDMAF      ;set seg base for buff (DMA)
jmp GETSEGTABLE ;return offset of Mem Desc Table
jmp GETIOBF      ;return I/O map byte (IOBYTE)
jmp SETIOBF      ;set I/O map byte (IOBYTE)

;***** INIT Entry Point, Differs for LDFIOS and BIOS, according to "Loader_Bios" value ****
;***** INIT Entry Point, Differs for LDFIOS and BIOS, according to "Loader_Bios" value ****

INIT:    ;print signon message and initialize hardware
        mov ax,cs          ;we entered with a JMPF so use
        mov ss,ax          ;CS: as the initial value of SS:.
        mov ds,ax          ;DS:,
        mov es,ax          ;and ES:
        ;use local stack during initialization
        mov sp,offset stkbase
        cld                ;set forward direction

        IF      not loader_bios
;-----
;|
;| This is a BIOS for the CPM.SYS file.
;| Setup all interrupt vectors in low
;| memory to address trap

        push ds            ;save the DS register
        mov IOBYTE,0        ;clear IOBYTE
        mov ax,0
        mov ds,ax
        mov es,ax          ;set ES and DS to zero
        ;setup interrupt 0 to address trap routine
        mov int0_offset,offset int_trap
        mov int0_segment,CS
        mov di,4
        mov si,0            ;then propagate
        mov cx,510          ;trap vector to
        rep movs ax,ax      ;all 255 interrupts
        ;BDOS offset to proper interrupt
        mov bdos_offset,bdos_ofst
        pop ds              ;restore the DS register

;| (additional CP/M-86 initialization)
;|
;-----ENDIF ;not loader_bios

        IF      loader_bios

```

```

;-----|
;|      ;This is a BIOS for the LOADER
;|      push ds          ;save data segment
;|      mov ax,0
;|      mov ds,ax        ;point to segment zero
;|      ;BDOS interrupt offset
;|      mov bdos_offset,bios_ofst
;|      mov bios_segment,CS ;bios interrupt segment
;|      (additional LOADER initialization)
;|      pop ds          ;restore data segment
;|
;-----|
;|      ENDIF    ;loader_bios
;
;-----|
;|      mov bx,offset signon
;|      call pmsg        ;print signon message
;|      mov cl,0          ;default to dr A: on coldstart
;|      jmp ccp          ;jump to cold start entry of CCP
;
;#BOOT: jmp ccp+6          ;direct entry to CCP at command level
;
;-----|
;|      IF      not loader_bios
;|
;-----|
;|      int_trap:
;|          cli           ;block interrupts
;|          mov ax,cs
;|          mov ds,ax        ;set our data segment
;|          mov bx,offset int_trp
;|          call pmsg
;|          hlt           ;hardstop
;|
;-----|
;|      ENDIF    ;not loader_bios
;
;***** CP/M Character I/O Interface Routines *****
;*
;*      console is USART (i8251A) on i8086 8612
;*          at ports D8/DA
;***** *****
;
;COVST:          ;console status
;    in al,csts
;    and al,2
;    jz const_ret
;    or al,255       ;return non-zero if rda
;const_ret:
;    ret            ;rcvr data available

```

```

CONIN:           ;console input
    call CONST
    jz CONIN
    in al,cdata
    and al,7fh      ;read data & remove parity bit
    ret

CONOUT:          ;console output
    in al,csts
    and al,1        ;get console status
    jz CONOUT
    mov al,cl
    out cdata,al    ;transmitter buffer is empty
    ret             ;then return data

LISTOUT:         ;list device output
    ret             ;not implemented

LISTST:          ;poll list status
    ret             ;not implemented

PUNCH:           ;write punch device
    ret             ;not implemented

READER:          ;return eor for now
    mov al,1ah
    ret

GETIOBF:         MOV AL,0      ;IOBYTE NOT IMPLEMENTED
    ret

SETIOBF:         ret          ;iobyte not implemented

; Routine to set and echo a console character
; and shift it to upper case

UCONECHO:
    call CONIN      ;get a console character
    push ax
    mov cl,al      ;save and
    call CONOUT
    pop ax
    cmp al,'a'
    jb uret       ;less than 'a' is ok
    cmp al,'z'
    ja uret       ;greater than 'z' is ok

```

```

        sub al,'a'-'A' ;else shift to caps
uret:    ret

ptsg:    mov al,[BX]      ;get next char from message
          test al,al
          jz return      ;if zero return
          mov CL,AL
          call CONOUT    ;print it
          inc BX
          jmps ptsg      ;next character and loop

;*****
;*
;*           Disk Input/Output Routines
;*
;*****
SELDSK:   ;select disk given by register CL
ndisks equ 2 ;number of disks (up to 16)
          mov disk,cl ;save disk number
          mov bx,0000h ;ready for error return
          cmp cl,ndisks ;n beyond max disks?
          jnb return    ;return if so
          mov ch,0       ;double(n)
          mov bx,cx     ;bx = n
          mov cl,4       ;ready for *16
          shr bx,cl     ;n = n * 16
          mov cx,offset ipbase
          add bx,cx     ;ipbase + n * 16
return:   ret          ;bx = .dpn

HOME:    ;move selected disk to home position (Track 0)
          mov io_com,homcom
          mov trk,0
          call execute
          ret

SETTRK:  ;set track address given by CL
          mov trk,CL
          ret

SETSEC:  ;set sector number given by cl
          mov sect,CL
          ret

SECTRAN: ;translate sector CX using table at [DX]
          mov ch,0
          mov bx,cx
          add bx,dx      ;add sector to tran table address
          mov bl,[bx]      ;get logical sector

```

```

        ret

SETDMA: ;set DMA offset given by CX
        mov dma_adr,CX
        ret

SETDMA_B: ;set DMA segment given by CX
        mov dma_seg,CX
        ret
;

GETSEG_T: ;return address of physical memory table
        mov bx,offset seg_table
        ret

;***** All disk I/O parameters are setup: *****
;*      DISK      is disk number      (SELDSK) *
;*      TRK       is track number    (SETTRK) *
;*      SECT      is sector number   (SETSEC) *
;*      DMA_ADDR  is the DMA lsb offset   *
;*      READ reads the selected sector to the DMA *
;*      address, and WRITE writes the data from   *
;*      the DMA address to the selected sector   *
;*      *
;***** All disk I/O parameters are setup: *****

READ:
        mov cl,4
        mov al,disk      ;combine disk selection
        sal al,cl        ;with opcode
        or al,ricode
        mov io_com,al    ;create iopb
        jmps execute

WRITE:
        mov cl,4
        mov al,disk
        sal al,cl
        or al,wicode    ;create iopb for write
        mov io_com,al

EXECUTE:

outer_retry:
        mov rtry_cnt,max_retries

retry:
        in al,rtype      ;clear controller
        in al,rbyte
        call sendcom

```

```

idle:    in al,istat      ;wait for completion
         and al,4          ;ready
         jz idle

;       check i.o. completion ok
;       in al,rtype
;       00 unlinked i/o complete           01 linked i/o comp
;       12 disk status changed          11 (not used)
;       must be a 00 in al
;       test al,100      ;ready status change?
;       JNZ WREADY
;       OR AL,2
;       jnz werror      ;some other error, retry

;       check i/o error bits
in al,rtype
rcf al,1
mov err_code,80h
jb wready      ;unit not ready
rcr al,1
mov err_code,al
and al,0feh    ;any other errors?
jnz werror

;       read or write is ok, al contains v
ret

wready: ;not ready, treat as an error for now
in al,rbyte      ;clear result byte
jmps trycount

werror: ;return hardware malfunction
trycount:
dec rtry_cnt
jnz retry
mov al,err_code
mov an,2
mov bx,ax      ;make error code 16 bits
mov bx,errtbl[BX]
call pmsg      ;print appropriate message
in al,cdata    ;flush usart receiver buffer
call uconecho   ;read upper case console character
cmp al,'C'
je wboot_i     ;cancel
cmp al,'R'
je outer_retry ;retry 10 more times
cmp al,'I'
je z_ret       ;ignore error
or al,255      ;set code for permanent error
z_ret: ret

```

```

wboot_l:           ;can't make it w/ a short leap
    jmp WBOOT

;*****
;*
;* sendcom sends the address of the iopb to
;*     the iSBc 201
;*
;*****
;***** sendcom: *****

sendcom:
    MOV CL,4
    MOV AX,DMA_SEG
    SAL AX,CL
    ADD AX,DMA_ADR
    MOV IO_ADR,AX
    MOV CL,4
    MOV AX,CS
    SAL AX,CL
    ADD AX,OFFSET CHANCMD ;ADD SEG & OFFSET FOR 201
    out ilow,al
    mov cl,8
    sar ax,cl
    out inigh,al
    ret

;*****
;*
;*          Data Areas
;*
;*****
data_offset      equ offset $


dseg
org      data_offset ;contiguous with code segment
IOBYTE  db      0
disk    db      0 ;disk number
chancmd db      80h ;iopb channel word
io_com   db      0
nsec    db      1 ;number sectors to xfer
trk     db      0
sect    db      0 ;start sector
IO_ADR  DW      0000H ;PHYS ADDR FOR SBC201 USE
dma_adr dw      0080h ;DMA adr (default)
dma_seg dw      0 ;DMA Base Segment

HOM_COM EQU 3
RDCODE  EQU 4

```

```

ERR_CODE DB 26H
VRCODE EQU 6

        IF      loader_bios
;-----+
;|
signon db      cr,lf,cr,lf
        db      'CP/M-86 Version 1.2',cr,lf,0
;|
;-----+
        ENDIF ;loader_bios

        IF      not loader_bios
;-----+
;|
signon db      cr,lf,cr,lf
        db      'System Generated 04/28/81'
        db      cr,lf,0
;|
;-----+
        ENDIF ;not loader_bios

int_trp db      cr,lf
        db      'Interrupt Trap Halt'
        db      cr,lf,0

errtbl dw er0,er1,er2,er3
dw er4,er5,er6,er7
dw er8,er9,erA,erB
dw erC,erD,erE,erF
dw er10,er20,er40,er80

er0    db      cr,lf,'Null Error ??',0
er1    db      cr,lf,'Deleted Record :',0
er2    db      cr,lf,'CRC Error :',0
er3    db      cr,lf,'Data Overrun-Underrun :',0
er4    db      cr,lf,'Seek Error :',0
er5    equ er0
er6    equ er0
er7    equ er0
er8    db      cr,lf,'Address Error :',0
er9    db      cr,lf,'Write Protect :',0
erA    db      cr,lf,'ID CRC Error :',0
erB    db      cr,lf,'Write Error :',0
erC    db      cr,lf,'Sector Not Found :',0
erD    equ er0
erE    db      cr,lf,'No Address Mark :',0
erF    db      cr,lf,'Data Mark Error :',0
er10   equ er3
er20   equ er9
er40   equ erB

```

```

er80    db cr,if,'Drive Not Ready :',e
rtry_cnt db 0      ;disk error retry counter

;       System Memory Segment Table

segtable db 1      ;1 segments
        dw tpa_seg      ;1st seg starts after BIOS
        dw tpa_len       ;and extends to 28000
include singles.lib ;read in disk definitions

loc_stk rw 32      ;local stack for initialization
stkbase equ offset $

lastoff equ offset $
tpa_seg equ (lastoff+04000n+15) / 16
tpa_len equ 0F00h - tpa_seg
        db 0      ;fill last address for GENCMD

;*****#
;**          Dummy Data Section
;**#
;*****#
dseg    E      ;absolute low memory
org     0      ;(interrupt vectors)
int0_offset    rw    1
int0_segment   rw    1
;      pad to system call vector
        rw    2*(bdos_int-1)

bios_offset      rw    1
bdos_segment    rw    1
END

rtry_cnt db 0      ;disk error retry counter

```

## APPENDIX F

title "Customized Basic I/O System"

```
;*****  
;*  
;* This Customized BIOS adapts CP/M-86 to *  
;* the following hardware configuration *  
;* Processor: iSBC 8612 *  
;* Controller: iSBC 202 *  
;* Memory model: 8080 *  
;*  
;* Programmer: M.B. Candalor *  
;* Revisions : *  
;*  
;*****  
  
true equ -1  
false equ not true  
cr equ 0dh ;carriage return  
lf equ 0an ;line feed  
max_retries equ 10 ;for disk i/o, before perr error  
  
;*****  
;*  
;* Loader_bios is true if assembling the *  
;* LOADER BIOS, otherwise BIOS is for the *  
;* CPM.SYS file.  
;*  
;*****  
  
loader_bios EQU TRUE  
bios_int equ 224 ;reserved BDOS interrupt  
  
IF not loader_bios  
----  
;  
bios_code equ 2500h  
ccp_offset equ 0000h  
bios_ofst equ 0B06h ;BDOS entry point  
;  
----  
ENDIF ;not loader_bios  
  
IF loader_bios  
----  
;  
bios_code equ 1200h ;start of LDBIOS  
ccp_offset equ 0003h ;base of CPMILOADER  
bios_ofst equ 0426h ;stripped BDOS entry
```

```

;-----[ENDIF ;loader_bios]-----;

csts    equ 0dah      ;I8251 status port
cdata   equ 0d8h      ;          data
;
;
;***** INTEL iSBC 202 Disk Controller Ports *****
;*
;* INTEL iSBC 202 Disk Controller Ports
;*
;***** ***** ***** ***** ***** ***** ***** *****

base    equ 078h
rtype   equ base+1
rbyte   equ base+3
reset   equ base+7

dstat   equ base
ilow    equ base+1
inigh   equ base+2

        cseg
        org  ccpoffset
ccp:
        org  bios_code

;***** BIOS Jump Vector for Individual Routines *****
;*
;* BIOS Jump Vector for Individual Routines
;*
;***** ***** ***** ***** ***** ***** ***** *****

jmp INIT           ;Enter from BOOT ROM or LOADER
jmp WBOOT          ;Arrive here from BDOS call 0
jmp CONST          ;return console keyboard status
jmp CONIN          ;return console keyboard char
jmp CONOUT         ;write char to console device
jmp LISTOUT        ;write character to list device
jmp PUNCH          ;write character to punch device
jmp READER         ;return char from reader device
jmp HOME           ;move to trk 00 on cur sel drive
jmp SELDSK         ;select disk for next rd/write
jmp SETTRK         ;set track for next rd/write
jmp SETSEC         ;set sector for next rd/write
jmp SETDMA         ;set offset for user buff (DMA)
jmp READ            ;read a 128 byte sector
jmp WRITE           ;write a 128 byte sector
jmp LISTST          ;return list status

```

```

jnp SECTRAN      ;xlate logical->physical sector
jnp SETDMAAB    ;set sec base for buff (DMA)
jnp GETSEGTT    ;return offset of Mem Desc Table
jnp GETIOBF      ;return I/O map byte (IOBYTE)
jnp SETIOBF      ;set I/O map byte (IOBYTE)

;*****
;* INIT Entry Point, Differs for LDBIOS and *
;* BIOS, according to "Loader_Bios" value   *
;*                                              *
;*****


INIT:    ;print signon message and initialize hardware
        mov ax,cs          ;we entered with a JMPF so use
        mov ss,ax          ;CS: as the initial value of SS:
        mov ds,ax          ;DS:
        mov es,ax          ;and ES:
        ;use local stack during initialization
        mov sp,offset stkbase
        cld                ;set forward direction

        IF      not loader_bios
        -----
        ;|
        ; This is a BIOS for the CPM.SYS file.
        ; Setup all interrupt vectors in low
        ; memory to address trap

        push ds            ;save the DS register
        mov IOBYTE,0        ;clear IOBYTE
        mov ax,0
        mov ds,ax
        mov es,ax          ;set ES and DS to zero
        ;setup interrupt 0 to address trap routine
        mov int0_offset,offset int_trap
        mov int0_segment,CS
        mov di,4
        mov si,0            ;then propagate
        mov cx,512          ;trap vector to
        rep movs ax,ax      ;all 256 interrupts
        ;BDOS offset to proper interrupt
        mov bdos_offset,bdos_ofst
        pop ds              ;restore the DS register

        ;(additional CP/M-SG initialization)
        ;|
        ;-----
        ENDIF    ;not loader_bios

        IF      loader_bios

```

```

;-----!
;|
;| ;This is a BIOS for the LOADER
;| push ds           ;save data segment
;| mov ax,4
;| mov ds,ax         ;point to segment zero
;| ;BDOS interrupt offset
;| mov bios_offset,bios_ofst
;| mov bdos_segment,CS ;bdos interrupt segment
;| ;(additional LOADER initialization)
;| pop ds           ;restore data segment
;|
;-----!
;|
;| ENDIF    ;loader_bios
;
;| mov bx,offset signon
;| call pmsg          ;print signon message
;| mov cl,0            ;default to dr A: on coldstart
;| jmp ccp            ;jump to cold start entry of CCP
;
;WBOOT: jmp ccp+6          ;direct entry to CCP at command level
;
;| IF      not loader_bios
;-----!
;|
;| int_trap:
;|     cli             ;block interrupts
;|     mov ax,cs
;|     mov ds,ax         ;get our data segment
;|     mov bx,offset int_trp
;|     call pmsg
;|     hlt             ;hardstop
;|
;-----!
;|
;| ENDIF    ;not loader_bios
;
;***** CP/M Character I/O Interface Routines *****
;*
;*      console is USART (i8251A) on i86C 8612
;*      at ports D8/DA
;*****
;
CONST:      ;console status
    in al,csts
    and al,2
    jz const_ret
    or al,255      ;return non-zero if raa
const_ret:
    ret            ;rcvr data available

```

```

CONIN:           ;console input
    call CONST
    jz CONIN
    in al,cdata
    and al,7fh      ;read data & remove parity bit
    ret

CONOUT:          ;console output
    in al,csts
    ani al,1        ;get console status
    jz CONOUT
    mov al,cl
    out cdata,al    ;transmitter buffer is empty
    ret             ;then return data

LISTOUT:         ;list device output
    ;not implemented
    ret

LISTST:          ;poll list status
    ;not implemented
    ret

PUNCH:           ;write punch device
    ;not implemented
    ret

READER:          ;return eof for now
    mov al,lan
    ret

GETIOPF:          MOV AL,0      ;IOBYTE NOT IMPLEMENTED
    ret

SETIOPF:          ret          ;iobyte not implemented

; Routine to get and echo a console character
; and shift it to upper case

uconecho:
    call CONIN      ;get a console character
    push ax
    mov cl,al
    call CONOUT
    pop ax
    cmp al,'a'
    jb uret
    cmp al,'z'
    ja uret
    ;less than 'a' is ok
    ;greater than 'z' is ok

```

```

        sub al,'a'-'A' ;else shift to caps
uret:
        ret
pmsg:
        mov al,[BX]      ;get next char from message
        test al,al
        jz return       ;if zero return
        mov CL,AL
        call CONOUT     ;print it
        inc BX
        jnps pmsg       ;next character and loop

;***** Disk Input/Output Routines *****
;*
;*
;*
SELDSK:    ;select disk given by register CL
ndisks equ 2 ;number of disks (up to 16)
        mov disk,CL ;save disk number
        mov bx,00000n ;ready for error return
        cmp cl,ndisks ;n beyond max disks?
        jnb return ;return if so
        mov cn,0 ;double(n)
        mov bx,cx ;bx = n
        mov cl,4 ;ready for *16
        shr bx,cl ;n = n * 16
        mov cx,offset ipbase
        add bx,cx ;ipbase + n * 16
return: ret ;bx = .dpn

HOME:   ;move selected disk to home position (Track 0)
        mov io_com,notcom
        mov trk,0
        call execute
        ret

SETTRK: ;set track address given by CL
        mov trk,CL
        ret

SETSEC: ;set sector number given by cl
        mov sect,CL
        ret

SECTRAN: ;translate sector CX using table at [DX]
        mov cn,0
        mov bx,cx
        TEST DX,00           ;IS THERE A SKEW?
        JZ NO_SKW            ;IF NOT, RET

```

```

        add bx,dx      ;add sector to tran table address
        mov bl,[bx]      ;set logical sector
        ret

NO_SKW:
        ADD BX,1
        RET

SETDMA: ;set DMA offset given by CX
        mov dma_adr,CX
        ret

SETDMAB: ;set DMA segment given by CX
        mov dma_seg,CX
        ret
;

GETSEGت: ;return address of physical memory table
        mov bx,offset seg_table
        ret

;***** All disk I/O parameters are setup: *****
;* DISK      is disk number      (SELDISK) *
;* TRK       is track number    (SETTRK)  *
;* SECT      is sector number   (SETSEC)  *
;* IO_ADDR   IS THE PHYS ADDR FOR DMA   *
;* DMA_ADDR  is the DMA lsc offset   *
;* READ reads the selected sector to the DMA*
;* address, and WRITE writes the data from   *
;* the DMA address to the selected sector   *
;*                                         *
;*****                                         *****

READ:
        mov cl,4
        mov al,disk      ;combine disk selection
        sal al,cl        ;with opcode
        or al,rdcode
        mov io_com,al    ;create iopb
        jmps execute

WRITE:
        mov cl,4
        mov al,disk
        sal al,cl
        or al,wrcode    ;create iopb for write
        mov io_com,al

EXECUTE:
outer_retry:

```

```

        mov rtry_cnt,max_retries

retry:
        in al,rtype      ;clear controller
        in al,rbyte      ;
        call sendcom

idle:   in al,istat      ;wait for completion
        and al,4          ;ready
        jz idle

;       check i.o. completion ok
        in al,rtype
        ;00 unlinked i/o complete           01 linked i/o comp
        ;10 disk status changed           11 (not used)
        ;must be a 00 in al
        test al,10b      ;ready status change?
        JNZ WREADY
        OR AL,0
        JNZ werror      ;some other error, retry

;       check i/o error bits
        in al,rbyte
        rcl al,1
        mov err_code,80h
        Jb wready      ;unit not ready
        rcr al,1
        mov err_code,al
        and al,0feh      ;any other errors?
        jnz werror

;       read or write is ok, al contains 0
        ret

wready: ;not ready, treat as an error for now
        in al,rtype      ;clear result byte
        jmps trycount

werror: ;return hardware malfunction
trycount:
        dec rtry_cnt
        jnz retry
        mov al,err_code
        mov ah,0
        mov bx,ax      ;make error code 16 bits
        mov bx,errtbl[BX]
        call pmse      ;print appropriate message
        in al,cdata      ;flush usart receiver buffer
        call uconecho    ;read upper case console character
        cmp al,'C'
        je wboot_1      ;cancel

```

```

        cmp al,'R'
        je outer_retry ;retry 10 more times
        cmp al,'I'
        je z_ret         ;ignore error
        or al,255       ;set code for permanent error
z_ret:  ret

wboot_l:           ;can't make it w/ a short leap
        jmp #BOOT

;***** sendcom sends the address of the iopb to
;**      the iSBC 202
;*****
;***** sendcom:                                     *****
MOV CL,4
MOV AX,DMA_SEG
SAL AX,CL
ADD AX,DMA_ADR
MOV IO_ADR,AX
MOV CL,4
MOV AX,CS
SAL AX,CL
ADD AX,OFFSET CHANCMD ;ADD SEG & OFFSET FOR 242
out ilow,al
mov cl,8
sar ax,~1
out inien,al
ret

;***** Data Areas
;*****
;***** data_offset   equ offset S
;***** dseg
;***** org  data_offset    ;contiguous with code segment
IOBYTE db 0
disk  db 2      ;disk number
chancmd db $0n   ;iopb channel word
io_com db 0
usec   db 1      ;number sectors to xfer
trk    db 0
sect   db 0      ;start sector

```

```

IO_ADR DW      0022H ;PHYS ADDR FOR SEC202 USE
dma_adr dw      00B0n ;DMA adr (default)
dma_seg dw      0       ;DMA Base Segment

HOM_COM EQU 3
RDCODE EQU 4
ERR_CODE DB 22H
WRCODE EQU 5

        IF      loader_bios
;-----|
;|
signon db      cr,lf,cr,lf
        db      'CP/M-86 Version 1.0',cr,lf,%
;|
;-----|
        ENDIF    ;loader_bios

        IF      not loader_bios
;-----|
;|
signon db      cr,lf,cr,lf
        db      'System Generated 05/25/81'
        db      cr,lf,%
;|
;-----|
        ENDIF    ;not loader_bios

int_trp 1b      cr,if
        db      'Interrupt Trap Halt'
        db      cr,lf,%

errtbl dw  er0,er1,er2,er3
        dw  er4,er5,er6,er7
        dw  er8,er9,erA,erB
        dw  erC,erD,erE,erF
        dw  er10,er20,er40,er80

er0      db      cr,lf,'Null Error ??',%
er1      db      cr,lf,'Deleted Record :',%
er2      db      cr,lf,'CRC Error :',%
er3      db      cr,lf,'Data Overrun-Underrun :',%
er4      db      cr,lf,'Seek Error :',%
er5      equ  er0
er6      equ  er0
er7      equ  er0
er8      db      cr,lf,'Address Error :',%
er9      db      cr,lf,'Write Protect :',%
erA      db      cr,lf,'ID CRC Error :',%
erB      db      cr,lf,'Write Error :',%
erC      db      cr,lf,'Sector Not Found :',%

```



## APPENDIX C

```
;  
; ROM bootstrap for CP/M-86 on an iSBC86/12  
; with the  
; iSBC 221 & 222 Floppy Disk Controllers  
;  
;  
; Copyright (C) 1982,1981  
; Digital Research, Inc.  
; Box 579, Pacific Grove  
; California, 93950  
;  
;*****  
;* This is the BOOT ROM which is resident *  
;* in the 957 monitor. To execute the boot *  
;* the monitor must be brought on-line and *  
;* then control passed by the command *  
;* "effd4:2". First, the ROM moves *  
;* a copy of its data area to RAM at loca- *  
;* tion 00020H, then initializes the segment *  
;* registers and the stack pointer. The *  
;* various peripheral interface chips on the *  
;* SBC 86/12 are initialized. The 8251 *  
;* serial interface is configured for a 9600*  
;* baud asynchronous terminal, and the in- *  
;* terrupt controller is setup for inter- *  
;* rupts 10H-17H (vectors at 00040H-0005FH) *  
;* and edge-triggered auto-EOI (end of in- *  
;* terrupt) mode with all interrupt levels *  
;* masked-off. Next, the 221-222 Diskette *  
;* controller is initialized, and track 0 *  
;* sector 1 is read to determine the target *  
;* paragraph address for LOADER. Finally, *  
;* the LOADER on track 0 sectors 2-26 and *  
;* track 1 sectors 1-26 is read into the *  
;* target address. Control then transfers *  
;* to LOADER. ROM  
;* 0 contains the even memory locations, and  
;* ROM 1 contains the odd addresses. BOOT *  
;* ROM uses RAM between 00000H and 00FFFH *  
;* (absolute) for a scratch area, along with  
;* the sector 1 buffer.  
;*****  
;  
cr          equ      15  
lf          equ      10  
;  
;          disk ports and commands  
;
```

```

base          equ      278h
rtype         equ      base+1
rbyte        equ      base+3
reset         equ      base+7
;
istat         equ      base
ilow          equ      base+1
inigh         equ      base+2
;
;actual console baud rate
baud_rate    equ      9600
;value for s253 baud counter
baud         equ      768/(baud_rate/100)
;
csts          equ      0DAn    ;15251 status port
cdata         equ      0D8h    ;"      data port
;
tcn0          equ      0D0h    ;8253 PIC channel 0 port
tcn1          equ      tcn0+2 ;ch 1 port
tcn2          equ      tcn0+4 ;ch 2 port
tcnd          equ      tcn0+6 ;8253 command port
;
icp1          equ      2C2h    ;8259a port 0
icp2          equ      0C2h    ;8259a port 1
;
secsec        equ      2c8h    ;offset for track 1
;
ROMSEG        EQU      0F1D4H
;
;
;
;
cseg          romseg
;
;First, move our data area into RAM at 0000:0200
;
    mov ax,cs
    mov ds,ax      ;point DS to CS for source
    mov SI,drombegin ;start of data
    mov DI,offset ram_start ;offset of destination
    mov ax,0
    mov es,ax      ;destination segment is 0000
    mov CX,data_length ;how much to move in bytes
        rep movs al,al      ;move out of eprom a byte
                                ;at a time
;
    mov ax,0
    mov ds,ax      ;data segment now in RAM
    mov SS,ax
    mov SP,stack_offset ;Initialize stack segment/
                        ;pointer
    cld           ;clear the direction flag

```

```

;
;Setup the 8259 Programmable Interrupt Controller
;
    mov al,13h
    out icp1,al      ;8259a ICW 1  8086 mode
    mov al,10h
    out icp2,al      ;8259a ICW 2  vector 0 40-5F
    mov al,1Fh
    out icp2,al      ;8259a ICW 4  auto EOI master
    mov al,0FFh
    out icp2,al      ;8259a OCW 1  mask all levels off
;
;Reset and initialize the 221/222 Diskette Interface
;
restart:           ;also come back here on fatal errors
    in al,rtype      ;clear status type register
    in al,rbyte      ;clear status register
    out reset,al      ;reset diskette system
homer:   mov BX,offset home
          CALL execute  ;home drive 0
;
    mov bx,OFFSET sector1  ;offset for first sector DMA
    mov ax,bx      ;enter in packet
    mov bx,offset read0+5  ;"
    mov [bx],al      ;"
    inc bx
    mov [bx],an      ;packet now complete
    mov bx,offset read0  ;packet location
    call execute      ;send packet
;
    mov es,abs        ;segment loc for LOADER
    mov ax,es        ;must translate to 16 bit abs
    mov cl,04        ;addr for diskette controller
    sal ax,cl
    mov bx,offset read1+5
    mov [bx],al      ;enter in packet
    inc bx
    mov [bx],an
    mov bx,offset read1
    call execute      ;read track 0
;
    mov cl,04
    mov ax,es
    add ax,secsec
    sal ax,cl
    mov bx,offset read2+5
    mov [bx],al
    inc bx
    mov [bx],an
    mov bx,offset read2
;
```

```

        call execute          ;read track 2
;
        mov leap_segment.es
;
        ;setup far jump vector
        mov leap_offset,&
;
;
;
;
;
        enter LOADER
        jmpf dword ptr leap_offset
;
pmsz:
        mov cl,[BX]
        test cl,cl
        jz return
        call conout
        inc BX
        jmp pmsz
;
conout:
        in al,csts
        test al,1
        jz conout
        mov al,cl
        out cdata,al
        ret
;
conin:
        in al,csts
        test al,2
        jz conin
        in al,cdata
        and al,7Fn
        ret
;
;
;
execute:
retry:           ;ret " if drive not ready
                ;clear controller
                ;      "
idle:            in al,dstat          ;system status
                and al,4
                jz idle           ;system awaiting interrupt
                in al,rtype
                test al,2         ;check drive status
                jz intcmp
                jmp RETRY        ;I/O NOT COMPLETE, TRY AGAIN

```

```

intemp: in al,rbyte           ;io is complete set status
        rcl al,1
        jae iocmp
        RCR AL,1
        jmp fatal
iocmp: rcr al,1               ;restore
        and al,EFEN             ;any errors ?
        jz return
        JMP RETRY

;
;

;
;

fatal:          ; fatal error
        mov cl,0
FTEST: RCR AL,1
        inc cl
        TEST AL,Z1
        jz ftest
        mov al,cl
        mov ah,0
        ADD AX,AX
        mov bx,ax                 ;make 16 bits
        mov bx,errtbl[EX]
;       print appropriate error message
        call pmsg
        call conin                ;wait for key strike
        pop ax                    ;discard unused item
        jmp restart               ;then start all over
;
return:         RET          ;return from EXECUTE
;
;
;
;

sendcom:        ;routine to send a command string to ZY1/271
        mov ax,bx
        out ilow,al
        mov cl,0B
        sar ax,cl
        out ihigh,al              ;packet adir
        ret

;
;
;
;

        Image of data to be moved to RAM
;
;

drombegin equ offset $           ;
```

```

creadstrine    db      80h      ;parameter block icw
                db      4h      ;read function code for drive
                db      1      ;# sectors to read
                db      2      ;track #
                db      1      ;start with sector 1
                db      2      ;will contain lower byte addr
                db      2      ;"      upper
;
creadtrkf     db      80h
                db      4h      ;read multiple
                db      25     ;# sectors to read
                db      0      ;track #
                db      2      ;start with 2
                db      0      ;addr for track 0 goes here
                db      0
;
creatrkl      db      80h
                db      4h
                db      26     ;sectors
                db      1      ;track #
                db      1      ;start with sector 1
                db      0      ;addr lsb
                db      0      ;addr msb
;
chome0        db      82h
                DB      03H
                db      0
                db      0
                db      0
                db      0
                db      0
                db      2
;
cerrtol dw      offset er0
                dw      offset er1
                dw      offset er2
                dw      offset er3
                dw      offset er4
                dw      offset er5
                dw      offset er6
                dw      offset er7
;
Cer0    db      cr,lf,'Null Error ??',0
Cer1    db      cr,lf,'CRC Error',0
Cer2    db      cr,lf,'Seek Error',0
Cer3    db      cr,lf,'Address Error',0
Cer4    db      cr,lf,'Data Overrun-Underrun',0
Cer5    db      cr,lf,'Write Protect',0
Cer6    db      cr,lf,'Write Error',0
Cer7    db      cr,lf,'Drive Not Ready',0
;
dromend equ offset $

```

```

;
data_length      equ  dromend-drombegin
;
;      reserve space in RAM for data area
;      (no hex records generated here)
;
dseg    0
org    2200h
;
ram_start        equ   $ 
read0            rb    ?      ;read track 0 sector 1
read1            rb    ?      ;read T0 S2-26
read2            rb    ?      ;read T1 S1-26
home             rb    ?      ;home drive 0
errtol           rw    8
er0               rb    length cer0      ;16
er1               rb    length cer1
er2               rb    length cer2
er3               rb    length cer3
er4               rb    length cer4      ;14
er5               rb    length cer5      ;11
er6               rb    length cer6      ;15
er7               rb    length cer7      ;17
;
leap_offset       rw    1
leap_segment     rw    1
;
;
stack_offset      rw    32      ;local stack
equ    offset $;stack from here down
;
;      T0 S1 read in here
sector1          equ    offset $ 
;
Ty                rb    1
Len               rw    1
Abs               rw    1      ;ABS is all we care about
Min               rw    1
Max               rw    1
end

```

## APPENDIX D

### title '8286 Disk I/O Drivers'

```
;*****  
;**  
;** Basic Input/Output System (BIOS) for  
;** CP/M-86 Configured for iSBC 85/12 with  
;** the iSBC 274 Floppy Disk Controller  
;**  
;** (Note: this file contains both embedded  
;** tabs and blanks to minimize the list file  
;** width for printing purposes. You may wish  
;** to expand the blanks before performing  
;** major editing.)  
;*****  
  
; Copyright (C) 1980,1981  
; Digital Research, Inc.  
; Box 579, Pacific Grove  
; California, 93950  
  
;(Permission is hereby granted to use  
;or abstract the following program in  
;the implementation of CP/M, MP/M or  
;CP/NET for the 8086 or 8088 Micro-  
;processor)  
  
true equ -1  
false equ not true  
  
;*****  
;**  
;** Loader_bios is true if assembling the  
;** LOADER BIOS, otherwise BIOS is for the  
;** CPM.SYS file. Blc_list is true if we  
;** have a serial printer attached to ILC8538  
;** Bdos_int is interrupt used for earlier  
;** versions.  
;**  
;*****  
  
loader_bios equ false  
blc_list equ true  
bdos_int equ 224 ;reserved BDOS Interrupt  
  
IF not loader_bios  
-----  
;!  
-----
```

```

bios_code      equ 2500h
cpu_offset     equ 1000h
bios_ofst      equ 2B26h ;BIOS entry point
;|
;-----|
        ENDIF ;not loader_bios

        IF      loader_bios
;-----|
;|
sics_code      equ 1200h ;start of LDBIOS
cpu_offset     equ 2723h ;base of CPM1LOADER
bios_ofst      equ 0406h ;stripped BDOS entry
;|
;-----|
        ENDIF ;loader_bios

cssts          equ 0DAh ;i8251 status port
cdtata         equ 0D8h ;      "      data port

        IF      bic_list
;-----|
;|
lsts           equ 41h ;2651 No. 0 on BLC8538 status port
lata           equ 40h ;      "      "      "      data port
bic_reset      equ 60h ;reset selected USARTS on BLC8538
;|
;-----|
        ENDIF ;bic_list

;***** Intel iSBC 204 Disk Controller Ports *****
;*
;*      Intel iSBC 204 Disk Controller Ports      *
;*
;***** ***** ***** ***** ***** ***** ***** *****

base204         equ 0a00h      ;SBC204 assigned address

f1c_com         equ base204+0  ;8271 FDC out command
f1c_stat        equ base204+0  ;8271 in status
f1c_parm        equ base204+1  ;8271 out parameter
f1c_rslt        equ base204+1  ;8271 in result
f1c_rst         equ base204+2  ;8271 out reset
dmac_air        equ base204+4  ;8257 DMA base address out
dmac_cont       equ base204+5  ;8257 out control
dmac_scan       equ base204+6  ;8257 out scan control
dmac_sadr       equ base204+7  ;8257 out scan address
dmac_node       equ base204+8  ;8257 out mode
dmac_stat       equ base204+9  ;8257 in status
f1c_sel         equ base204+9  ;FDC select port (not used)
f1c_segment     equ base204+10 ;segment address register

```

```

reset_204      equ base204+15 ;reset entire interface
max_retries    equ 12          ;max retries on disk i/o
cr              equ 010          ;before perm error
lf              equ 0ah          ;carriage return
lf              equ 0dah         ;line feed

cseg
ore   ccpoffset
ccp:
org   bios_code

; *****
; * BIOS Jump Vector for Individual Routines *
; *
; *****

jmp INIT        ;Enter from BOOT ROM or LOADER
jmp WBOOT       ;Arrive here from BDOS call &
jmp CONST       ;return console keyboard status
jmp CONIN       ;return console keyboard char
jmp CONOUT      ;write cnar to console device
jmp LISTOUT     ;write character to list device
jmp PUNCH       ;write character to punch device
jmp READER      ;return cnar from reader device
jmp HOME        ;move to trk 00 on cur sel drive
jmp SELDSK      ;select disk for next rd/write
jmp SETTRK      ;set track for next rd/write
jmp SETSEC      ;set sector for next rd/write
jmp SETDMA      ;set offset for user buff (DMA)
jmp READ         ;read a 128 byte sector
jmp WRITE        ;write a 128 byte sector
jmp LISTST      ;return list status
jmp SECTRAN     ;xlate logical->physical sector
jmp SETDMAB     ;set seg base for buff (DMA)
jmp GETSEGT     ;return offset of Mem Desc Table
jmp GETIOBF     ;return I/O map byte (10BYTE)
jmp SETIOBF     ;set I/O map byte (10BYTE)

; *****
; * INIT Entry Point, Differs for LUBIOS and *
; * BIOS, according to "Loader_Bios" value *
; *
; *****

INIT:    ;print siemon message and initialize hardware
        mov ax,cs           ;we entered with a JMPF so use
        mov ss,ax           ; CS: as the initial value of SS:,
        mov ds,ax           ;                 DS:, DS:,

```

```

        mov es,ax      ; and ES:
;use local stack during initialization
        mov sp,offset stkbase
        cld          ;set forward direction

        IF      not loader_bios
;-----!
;| ; This is a BIOS for the CPM.SYS file.
;| ; Setup all interrupt vectors in low
;| ; memory to address trap

        push ds      ;save the DS register
        mov ax,0
        mov ds,ax
        mov es,ax      ;set ES and DS to zero
;setup interrupt 0 to address trap routine
        mov int0_offset,offset int_trap
        mov int0_segment,CS
        mov di,4
        mov si,0      ;then propagate
        mov cx,510    ;trap vector to
        rep movs ax,ax ;all 256 interrupts
;BDOS offset to proper interrupt
        mov bdos_offset,bdos_ofst
        pop ds      ;restore the DS register

;*****#
;*
;* National "ECL 8538" Channel 2 for a serial
;* 9600 baud printer - this board uses 8 Si232C
;* netics 2551 Usarts which have on-chip baud
;* rate generators.
;*
;*****#
;*                                     *
;*                                     *

        mov al,0FFh
        out bic_reset,al ;reset all usarts on 8538
        mov al,4En
        out ldata+2,al ;set usart 0 in async 8 bit mode
        mov al,3En
        out ldata+2,al ;set usart 0 to 9600 baud
        mov al,37h
        out ldata+3,al ;enable Tx/Rx, and set up RTS,ETR
;|
;-----!
;| ENDIF      ;not loader_bios

        IF      loader_bios
;-----!
;|

```



```

        jz CONIN          ;wait for RDA
        in al,cdata
        and al,7fh        ;read data and remove parity bit
        ret

CONOUT:      ;console output
        in al,csts
        and al,1          ;get console status
        jz CONOUT         ;wait for TPE
        mov al,cl
        out cdata,al      ;Transmitter Buffer Empty
        ret               ;then return data

LISTOUT:      ;list device output
        IF      blc_list
;-----|
;|           call LISTST
        jz LISTOUT        ;wait for printer not busy
        mov al,cl
        out ldata,al      ;send char to TI 812
;|           |
;-----|
        ENDIF  ;blc_list
        ret

LISTST:       ;pcii list status
        IF      blc_list
;-----|
;|           in al,lsts
        and al,81h        ;look at both TxRDY and CTR
        cmp al,81h
        jnz zero_ret      ;either false, printer is busy
        or al,255          ;both true, LPT is ready
;|           |
;-----|
        ENDIF  ;blc_list
        ret

PUNCH:  ;not implemented in this configuration
READER:
        mov al,lan
        ret               ;return EOF for now

GETIOBF:
        mov al,0          ;TTY: for consistency

```

```

        ret          ;IOBYTE not implemented

SETIOPBF:
        ret          ;iobyte not implemented

zero_ret:
        and al,0
        ret          ;return zero in AL and flags

; Routine to set and echo a console character
; and shift it to upper case

uconenecho:
        call CONIN    ;get a console character
        push ax
        mov cl,al    ;save and
        call CONOUT   ;echo to console
        pop ax
        cmp al,'a'
        jb uret     ;less than 'a' is ok
        cmp al,'z'
        ja uret     ;greater than 'z' is ok
        sub al,'a'-'A' ;else shift to caps
uret:
        ret

; utility subroutine to print messages

pmse:
        mov al,[BX]    ;set next char from message
        test al,al
        jz return     ;if zero return
        mov CL,AL
        call CONOUT   ;print it
        inc BX
        jmps pmse    ;next character and loop

;***** Disk Input/Output Routines *****
;*
;*
;*
SELDSK:      ;select disk given by register CL
        mov bx,0000h
        cmp cl,2
        jno return   ;this BIOS only supports 2 disks
        mov al,80h
        cmp cl,0
        jne sel1     ;drive 1 if not zero
        mov al,40h
        ;else drive is 0

```

```

sell:    mov sel_mask,al ;save drive select mask
           ;now, we need disk parameter address
        mov cx,0
        mov bx,cx      ;BX = word(CL)
        mov cl,4
        shr bx,cl      ;multiply drive code * 16
        ;create offset from Disk Parameter Base
        add bx,offset dp_base
return:
        ret

HOME:   ;move selected disk to home position (Track 0)
        mov trk,0       ;set disk i/o to track zero
        mov bx,offset hom_com
        call execute
        jz return      ;home drive and return if OK
        mov bx,offset bad_hom ;else print
        call pmsg      ;Home Error
        jmps nome      ;and retry

SETTRK: ;set track address given by CX
        mov trk,cl      ;we only use 8 bits of track address
        ret

SETSEC: ;set sector number given by CX
        mov sect,cl      ;we only use 8 bits of sector address
        ret

SECTRAN: ;translate sector CX using table at [DX]
        mov bx,cx
        add bx,dx      ;add sector to tran table address
        mov bl,[bx]      ;get logical sector
        ret

SETDMA: ;set DMA offset given by CX
        mov dma_adr,CX
        ret

SETDMAB: ;set DMA segment given by CX
        mov dma_seg,CX
        ret
;

GETSEGt: ;return address of physical memory table
        mov bx,offset seg_table
        ret

*****
;*
;* All disk I/O parameters are setup: the *
;* Read and Write entry points transfer one *
;* sector of 128 bytes to/from the current *
;
```

```

;# DMA address using the current disk drive #
;#
;***** ****
READ:
    mov al,12h      ;basic read sector command
    jmps r_w_common

WRITE:
    mov al,2ah      ;basic write sector command

r_w_common:
    mov bx,offset io_com ;point to command string
    mov byte ptr 1[BX],al ;put command into string
;    fall into execute and return

execute: ;execute command string.
; [BX] points to length,
;       followed by Command byte,
;       followed by length-1 parameter bytes

    mov last_com,BX ;save command address for retries
outer_retry:
;allow some retrying
    mov rtry_cnt,max_retries
retry:
    mov BX,last_com
    call send_com ;transmit command to i8271
;check status poll

    mov BX,last_com
    mov al,1[BX]    ;get command op code
    mov cx,0E00h    ;mask if it will be "int req"
    cmp al,2ch
    jb exec_poll  ;ok if it is an interrupt type..
    mov cx,8080h    ;else we use "not command busy"
    and al,0fn
    cmp al,0ch    ;unless there isn't
    mov al,0
    ja exec_exit  ;      any result
;poll for bits in CH,
;      toggled with bits in CL

exec_poll:
    in al,fdc_stat ;read status
    and al,cn
    xor al,cl      ; isolate what we want to poll
    jz exec_poll  ;and loop until it is done

;Operation complete,
in al,fdc_rslt ; see if result code indicates error
and al,1en

```

```

        jz exec_exit      ;no error, then exit
                           ;some type of error occurred . . .
        cmp al,12h
        je dr_nrdy      ;was it a not ready drive ?
                           ;no,
dr_rdy: ; then we just retry read or write
        dec rtry_cnt
        jnz retry        ; up to 16 times

;      retries do not recover from the
;      hard error

        mov ah,0
        mov bx,ax          ;make error code 16 bits
        mov bx,errtbl[BX]
        call pmsg           ;print appropriate message
        in al,cdata         ;flush usart receiver buffer
        call uconecho        ;read upper case console character
        cmp al,'C'
        je wboot_l          ;cancel
        cmp al,'Q'
        je outer_retry       ;retry 16 more times
        cmp al,'I'
        je z_ret             ;ignore error
        or al,255            ;set code for permanent error
exec_exit:
        ret

dr_nrdy:      ;here to wait for drive ready
        call test_ready
        jnz retry           ;if it's ready now we are done
        call test_ready
        jnz retry           ;if not ready twice in row.
        mov bx,offset nrdynmsg
        call pmsg ; Drive Not Ready"
nrdy01:
        call test_ready
        jz nrdy01           ;now loop until drive ready
        jmps retry           ;then go retry without decrement
zret:
        and al,0
        ret                 ;return with no error code

wboot_l:       ;can't make it w/ a short leap
        jmp WBOOT

;***** *****
;*
;* The 18271 requires a read status command *
;* to reset a drive-not-ready after the      *
;* drive becomes ready                      *
;
```

```

;*
;***** *****
test_ready:
    mov dh, 40h      ;proper mask if dr 1
    test sel_mask,80h
    jnz nr1y2
    mov dh, 04h      ;mask for dr 0 status bit
nr1y2:
    mov bx,offset rds_com
    call send_com
dr_poll:
    in al,fdc_stat   ;get status word
    test al,80h
    jnz dr_poll     ;wait for not command busy
    in al,fdc_rslt  ;get "special result"
    test al,dh      ;look at bit for this drive
    ret             ;return status of ready

;*****
;*
;* Send_com sends a command and parameters  *
;* to the i8271: BX addresses parameters.  *
;* The DMA controller is also initialized  *
;* if this is a read or write               *
;*
;***** *****
send_com:
    in al,fdc_stat
    test al,80h      ;insure command not busy
    jnz send_com    ;loop until ready

    ;see if we have to initialize for a DMA operation

    mov al,1[bx]    ;get command byte
    cmp al,12h
    jne write_maybe ;if not a read it could be write
    mov cl,40h
    jmps init_dma  ;is a read command, go set DMA
write_maybe:
    cmp al,0an
    jne dma_exit    ;leave DMA alone if not read or write
    mov cl,80h      ;we have write, not read
init_dma:
;we have a read or write operation, setup DMA controller
; (CL contains proper direction bit)
    mov al,04h
    out dmac_mode,al ;enable dmac
    mov al,00
    out dmac_cont,al ;send first byte to control port

```

```

        mov al,cl
        out dmac_cont,al ;load direction register
        mov ax,dma_adr
        out dmac_air,al ;send low byte of DMA
        mov al,ah
        out dmac_air,al ;send high byte
        mov ax,dma_seg
        out fdc_segment,al ;send low byte of segment address
        mov al,an
        out fdc_segment,al ;then high segment address
        ana_exit:
        mov cl,[BX]      ;get count
        inc BX
        mov al,[BX]      ;get command
        or al,sel_mask ;merge command and drive code
        out fdc_com,al ;send command byte
        parm_loop:
        dec cl
        jz exec_exit    ;no (more) parameters, return
        inc BX          ;point to (next) parameter
        parm_poli:
        in al,fdc_stat
        test al,20h      ;test "parameter register full" bit
        jnz parm_poli   ;idle until parm reg not full
        mov al,[BX]
        out fdc_parm,al ;send next parameter
        jmps parm_loop  ;go see if there are more parameters

;*****
;*
;*           Data Areas
;*
;******
data_offset    equ offset $  

        ds  

        org     data_offset      ;contiguous with code segment
        IF      loader_bios
;-----!  

;|  

signon db      cr,lf,cr,lf
        db      'CP/M-86 Version 2.2',cr,lf,%
;|  

;-----!  

        ENDIF  ;loader_bios
        IF      not loader_bios
;-----!  

;|  

signon db      cr,lf,cr,lf

```

```

        db      ' System Generated - 11 Jan 81',cr,lf,?

;-----  

        ENDIF ;not loader_bios  

bad_nom db      cr,lf,'None Error',cr,lf,  

int_trp db      cr,lf,'Interrupt Trap Halt',cr,lf,  

errtbl dw er0,er1,er2,er3  

        dw er4,er5,er6,er7  

        dw er8,er9,erA,erB  

        dw erC,erD,erE,erF  

er0      db cr,lf,'Null Error ??',0  

er1      equ er0  

er2      equ er2  

er3      equ er0  

er4      db cr,lf,'Clock Error :',0  

er5      db cr,lf,'Late DMA :',0  

er6      db cr,lf,'ID CRC Error :',0  

er7      db cr,lf,'Data CRC Error :',0  

er8      db cr,lf,'Drive Not Ready :',0  

er9      db cr,lf,'Write Protect :',0  

era      db cr,lf,'Trk 00 Not Found :',0  

erb      db cr,lf,'Write Fault :',0  

erc      db cr,lf,'Sector Not Found :',0  

erd      equ er2  

erre      equ er2  

erf      equ er0  

nriymsg equ erg  

rtry_cnt db 0      ;disk error retry counter  

last_com dw 0      ;address of last command string  

dma_adr dw 0      ;dma offset stored here  

dma_seg dw 0      ;dma segment stored here  

sel_mask db 42h ;select mask, 40h or 80h  

;      Various command strings for 18271  

io_com  db 3      ;length  

rd_wr    db 2      ;read/write function code  

trk     db 0      ;track #  

sect    db 0      ;sector #  

nom_com db 2,29h,0      ;none drive command  

rds_com db 1,2ca      ;read status command  

;      System Memory Segment Table  

segtable db 2      ;2 segments  

        dw tpa_see      ;1st see starts after BIOS

```

```

        dw tpa_len      ;and extends to 18FFF
        dw 2000h        ;second is 20000 - 
        iw 2000h        ;3FFF (128K)

        include singles.lib ;read in disk definitions

loc_stk rw 32 ;local stack for initialization
stkbase equ offset $

lastoff equ offset $ 
tpa_seg equ (lastoff+0400h+15) / 16
tpa_len equ 0800h - tpa_seg
        db 0      ;fill last address for GENCMD

;*****
;*
;*          Dummy Data Section
;*
;*****
dseg    0      ;absolute low memory
        org    0      ;(interrupt vectors)
int0_offset    rw      1
int0_segment   rw      1
;      pad to system call vector
        rw      2*(ddos_int-1)

ddos_offset    rw      1
ddos_segment   rw      1
        END

```

## LIST OF REFERENCES

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